

Ensuring Correct Distributed Writes to Delta Lake in Rust with Formal Verification

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ORGANIZED BY 😂 databricks

A bit about myself

https://about.houqp.me

- Currently core software lead at Neuralink
 - <u>https://neuralink.com/careers/</u>
- Formally tech-lead of Data and AI platform at Scribd
 - <u>https://www.scribd.com/about/engineering</u>
 - Delta-rs is a team effort between Christian (<u>@xianwill</u>), Mykhailo (<u>@mosyp</u>) and Tyler (<u>@rtyler</u>).
- Authors and contributors of many open-source projects:
 - o <u>Delta-rs</u>
 - Apache Arrow
 - Apache Airflow
 - o <u>ROAPI</u>
 - <u>KOReader</u>



- Native Delta lake implementation in pure Rust
 - With Python and Ruby bindings
- Previous Data+AI 2021 talk:

https://databricks.com/session_na21/growing _the-delta-ecosystem-to-rust-and-pythonwith-delta-rs





Agenda

- Delta Lake 101
- Distributed concurrent table writes in delta-rs
- Formal verification overview
- Intro to stateright



Delta Lake 101



What is a Delta lake anyway?

- Brings ACID to traditional data warehouse solution
- Single data source for both batching and streaming workload



Traditional data warehouse





Delta Lake's transaction log





Delta Lake's transaction log





Delta Lake's transaction log





















Safe concurrent write in S3



S3 limitation

S3 doesn't support PUT if absent operation





Two design directions

- Use DynamoDB as the log store
 - Less complexity, a lot easier to implement
- Use DynamoDB to implement put if absent for S3
 - Fully transparent and compatible with other Delta Lake readers



Let's use a distributed lock?





DynamoDB as a distributed lock

How hard could it possibly be?



What if one of the writers **crashed** while holding the lock?



How hard could it possibly be?



If one writer crashed while holding the lock, it will **eventually expire so another writer can acquire it**.



How hard could it possibly be?





How hard could it possibly be?



If one writer **paused** while holding the lock, it will **eventually expire so another writer will acquire it**.



How hard could it possibly be?

Time	Writer A	Writer B
Т0	Acquired lock	
T1	Paused	
T2	Lock expired	
Т3		Acquired expired lock
T4	Resumed	
Τ5	I have the lock!	I have the lock!

If one writer **paused** while holding the lock, it will **eventually expire so another writer will acquire it**.



- Write commit to a temp S3 location
 - Copy from s3://table/{uuid}.json to s3://table/1.json is idempotent
- Atomically acquire the lock with recorded S3 copy operation
- Expired lock needs to be repaired by another writer



Time	Writer A	Writer B
Т0	Acquired lock with "copy uuid_a.json to 1.json"	
T1	Paused	
T2	Lock expired	
Т3		Acquired expired lock with "copy uuid_b.json to 1.json"
T4		Repair lock by executing "copy uuid_a.json to 1.json"
Т5	Resumed	
Т6	copy uuid_a.json to 1.json	
Т7	Release the lock	
Т8		Release the lock
Т9		Retry acquiring the lock with "copy uuid_b.json to 2.json"

- Full design discussion available at <u>https://github.com/delta-io/delta-rs/discussions/89</u>
- Dynamodb based lock implemented from scratch in Rust: <u>https://crates.io/crates/dynamodb_lock</u>
 - Shout out to my ex-colleague Mykhailo Osypov (<u>@mosyp</u>)



What if we have more than two writers?

Time	Writer A	Writer B	Writer C	Writer D	Writer
Т0	Acquired lock				
T1	Paused				
T2	Lock expired				
Т3		Acquired expired lock			
T4		Repair lock by executing			
Т5	Resumed				
Т6	Do work				
Т7	Release the lock				
Т8		Release the lock			



Enter formal verification



Informal verifications are everywhere

- Static type checks
 - Assert desired properties by enforcing type constraints
- Unit tests
 - Assert desired properties by matching inputs with expected outputs



Formal verification approaches

Model checking

- Exhaustive state exploration
- Deductive reasoning
 - Ensure conformance of system specifications through mathematical proofs



Model checker state exploration example Possible state 1

Time	Writer A	Writer B
Т0	Acquired lock	
T1	Paused	
Т2	Lock expired	
Т3		Tried to acquire lock
T4		

Writer A:

- Paused
- Held an expired lock

Writer B:

• Acquired an expired lock



Model checker state exploration example Possible state 2

Time	Writer A	Writer B
Т0	Acquired lock	
T1	Paused	
Т2		Tried to acquire lock
Т3	Locked expired	
Т4		

Writer A:

- Paused
- Held an expired lock

Writer B:

• Failed to acquire lock



Model checker industrial use

• TLA+

- AWS S3
- Azure cosmos DB
- Intel multi-processor cache-coherence protocols
- See <u>https://lamport.azurewebsites.net/tla/industrial-use.html</u>



Shortcomings of TLA+

- Steep learning curve
- Rough tooling
- Separation of proof and implementation



Intro to stateright



- Model checker as a Rust library
 - Comparison with TLA+: <u>https://www.stateright.rs/comparison-with-tlaplus.html</u>
- Homepage: <u>https://www.stateright.rs/</u>



Specify systems as state machines

```
pub trait Model: Sized {
   type State;
   type Action;
   fn init_states(&self) -> Vec<Self::State>①;
   fn actions(&self, state: &Self::State, actions: &mut Vec<Self::Action>);
   fn next_state(
        &self,
        last_state: &Self::State,
        action: Self::Action
    ) -> Option<Self::State>;
    fn properties(&self) -> Vec<Property<Self>>① { ... }
}
```


Specify systems as state machines

 See delta-rs's full stateright proof at <u>https://github.com/delta-io/delta-rs/blob/main/proofs/src/main.rs</u> (about 500 lines of Rust code)

High level state machine based actor interface

```
pub trait Actor: Sized {
    type Msg: Clone + Debug + Eq + Hash;
    type State: Clone + Debug + PartialEq + Hash;
    fn on_start(&self, id: Id, o: &mut Out<Self>) -> Self::State;
    fn on_msg(
       &self,
       id: Id,
        state: &mut Cow<'_, Self::State>,
       src: Id,
       msg: Self::Msg,
       o: &mut Out<Self>
    ) { . . . }
   fn on_timeout(
       &self,
        id: Id,
        state: &mut Cow<'_, Self::State>,
       o: &mut Out<Self>
    ) { ... }
```

Is formal verification worth it? Yes!

- The exercise of formalization forces you to really think through the design step by step
 - We discovered a couple of design bugs during this process
- Automated model checking caught one correctness bug that all of us missed:

https://github.com/delta-io/delta-rs/pull/540#issue-1097376239

Stateright Explorer API Docs Book Crate Source

Status

- Model: AtomicRenameSys
- States: 287,395
- Unique States: 116,479
- Progress: Done

Properties

- 🛕 Counterexample found: Always no overwrite
- 🔽 Safety holds: Always no unexpected rename
- 🗹 Safety holds: Always not retry on successful rename
- 💟 Liveness holds: Eventually all writer clean shutdown
- V Liveness holds: Eventually all source objects are purged
- A Counterexample found: Eventually all renames are performed
- V Example found: Sometimes lock contention

Path of Actions

1. Pre-init

- 2. Init 0
- TryAcquireLock(0)
- 4. NewVersionObjectCheckExists(0)
- 5. TryAcquireLock(1)
- TryAcquireLock(1)
- TryAcquireLock(1)
- 8. TryAcquireLock(1)
- 9. RepairObjectCheckExists(1)
- RepairObjectCopy(1)
- 11. UpdateLockData(1)
- 12. NewVersionObjectCheckExists(1)
- 13. NewVersionObjectCopy(0)
- 14. NewVersionObjectCopy(1)

Next Action Choices

- OldVersionObjectDelete(0)
- OldVersionObjectDelete(1)
- TryAcquireLock(2)

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SUM

Current State Complete State? Compact? AtomicRenameState { writer_ctx: [WriterContext { state: NewVersionObjectCopied, lock_data: LockData { dst: "0", src: "writer_0", **}**, acquired_expired_lock: false, released expired lock: false, rename_err: None, target_version: 0, }, WriterContext { state: NewVersionObjectCopied, lock_data: LockData { dst: "0", src: "writer_1", }, acquired_expired_lock: false, released_expired_lock: false, rename_err: None, target_version: 0, }, WriterContext { state: Init, lock_data: LockData { dst: "". src: "". }, acquired_expired_lock: false, released_expired_lock: false, rename_err: None, target_version: 0, }, lock: Some(GlobalLock { data: LockData {

det · "A"

Building efficient, safe and correct systems

Combining Rust and Formal verification

- Rust compiler guarantees absence of unsafe memory access bugs
 - No more segfaults and race conditions
- Formal verification removes logical bugs*
- Model checking as a library (stateright) to keep proofs and implementations in sync

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Thank you

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